

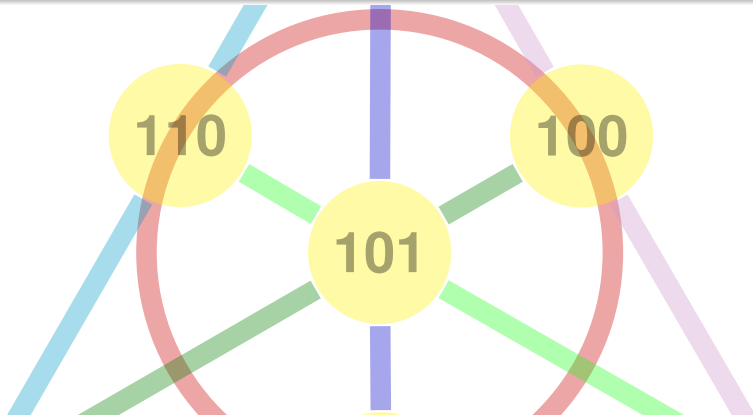
111

The Anti-Games Strike Back!

110

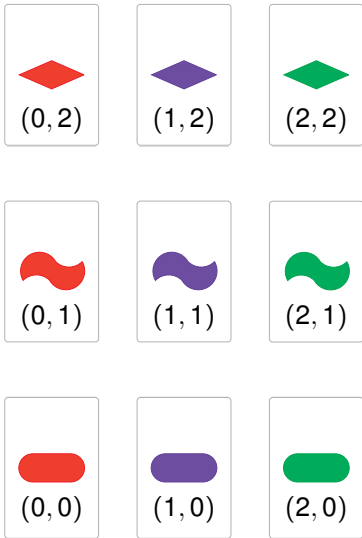
100

101



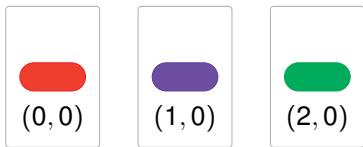
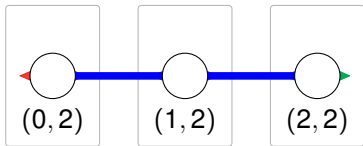
Let's play Anti-SET! ($AG(2)$)

First to get a set loses.



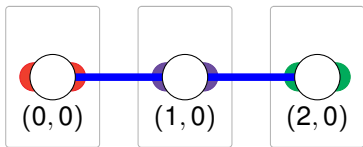
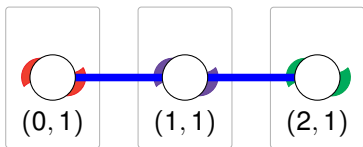
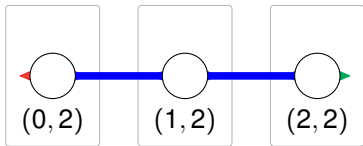
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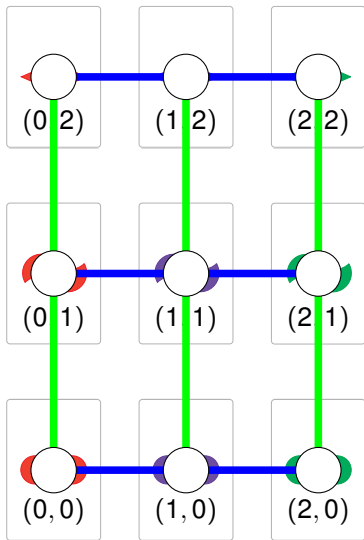
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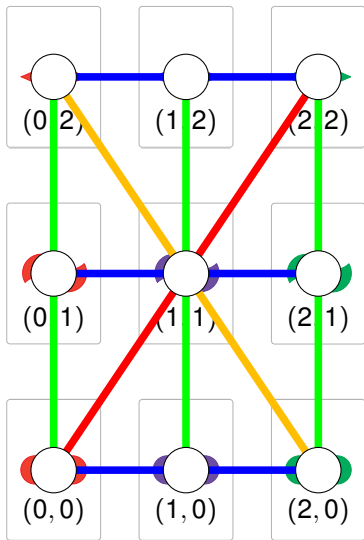
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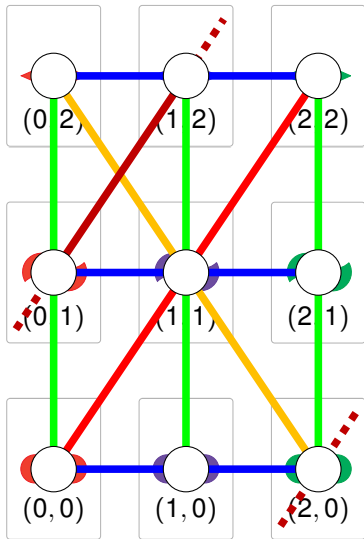
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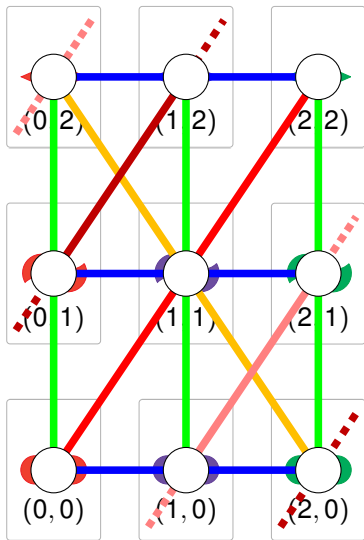
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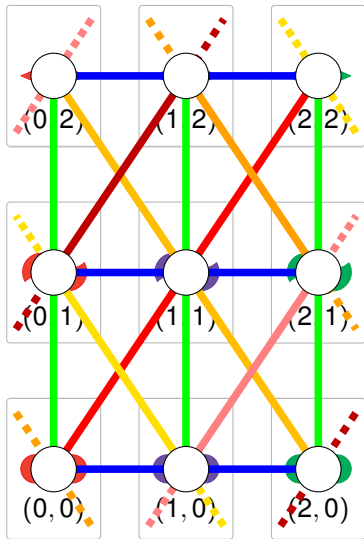
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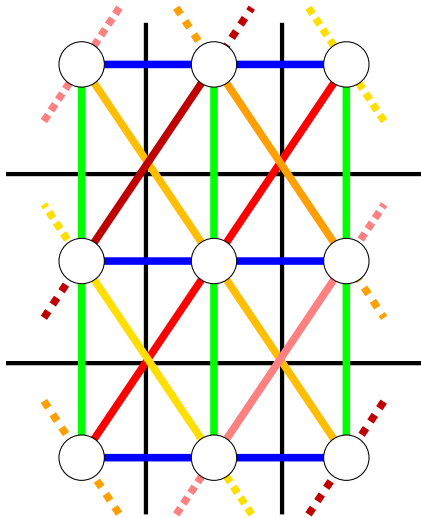
Let's play Anti-SET! ($AG(2)$)

First to get a set loses.



Let's play Anti-Tic-Tac-Toe!

First to get 3 in a row loses.



Anti-Tic-Tac-Toe

Anti-Tic-Tac-Toe is a game with two players:

Xavier (1st) and Olivia (2nd).

- The game is played on a “board” of points and lines.
- \mathcal{X} and \mathcal{O} alternate marking any unmarked point.
- A player loses immediately if they own a line.

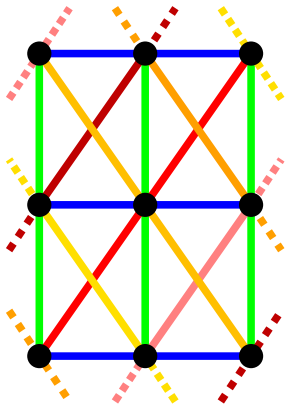
The fun depends on the “board”



Jacob Van Hook & Sophia Mancini
Mathfest 2017

What makes a good anti-tic-tac-toe board?

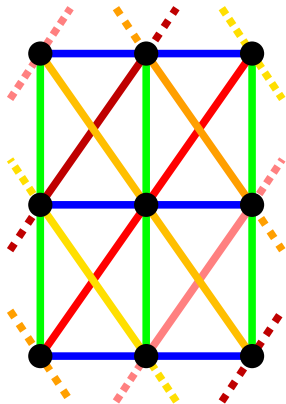
- 3 points per line.
- For every pair of points **there is a unique** line containing them.
- “Enough” points to avoid ties (> 3 points)



What makes a good anti-tic-tac-toe board?

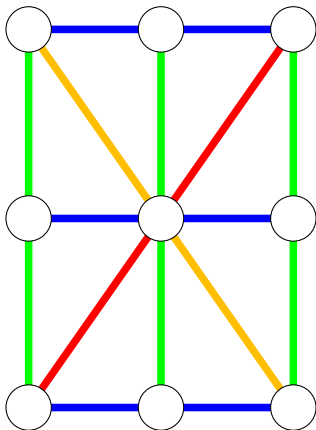
Steiner* Triple Systems!

- 3 points per line.
- For every pair of points **there is a unique** line containing them.
- “Enough” points to avoid ties (> 3 points)

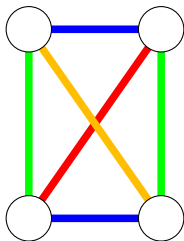


* Stigler's law

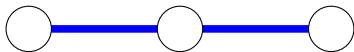
Steiner triple system or not?



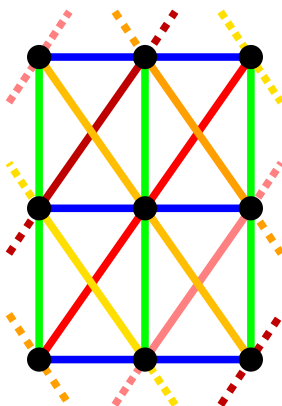
Steiner triple system or not?



Steiner triple system or not?



Last time ...



We showed that **Xavier wins** anti-tic-tac-toe
on **affine Steiner Triple Systems (AGs)**

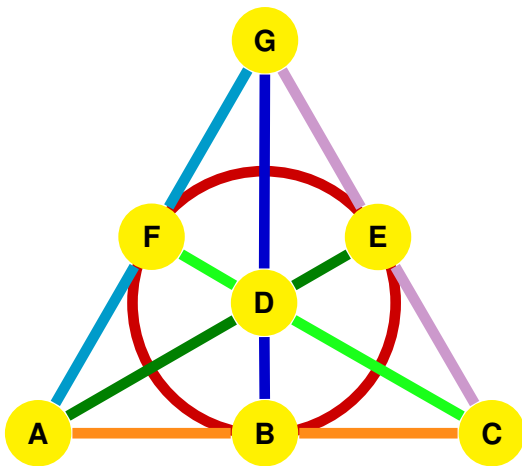
What other STSs are there?

What other STSs are there?



Projective Geometry

Another kind of Steiner Triple System



Projective Geometry $PG(3)$

Points: Nonzero binary points (p_1, p_2, p_3)

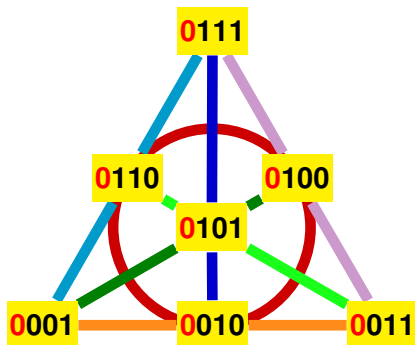
Lines: $\ell + m \equiv n \pmod{2}$



Projective Geometry $PG(4)$

Points: Nonzero binary points (p_1, p_2, p_3, p_4)

Lines: $\ell + m \equiv n \pmod{2}$

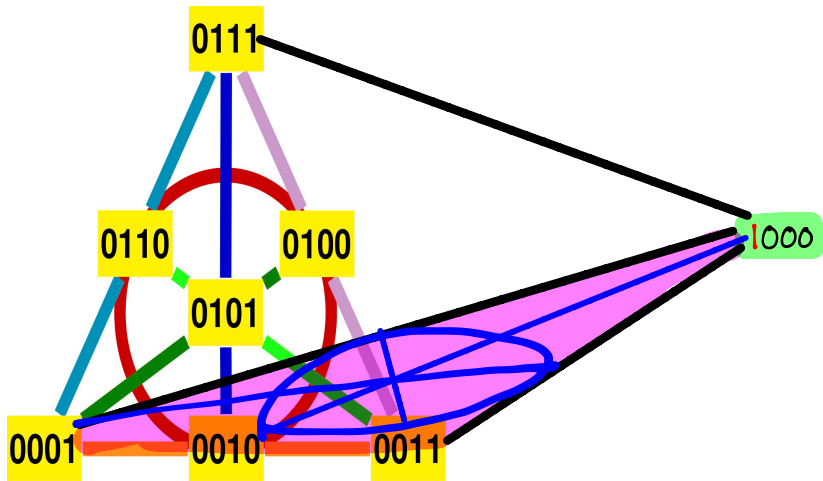


1000

Projective Geometry $PG(4)$

Points: Nonzero binary points (p_1, p_2, p_3, p_4)

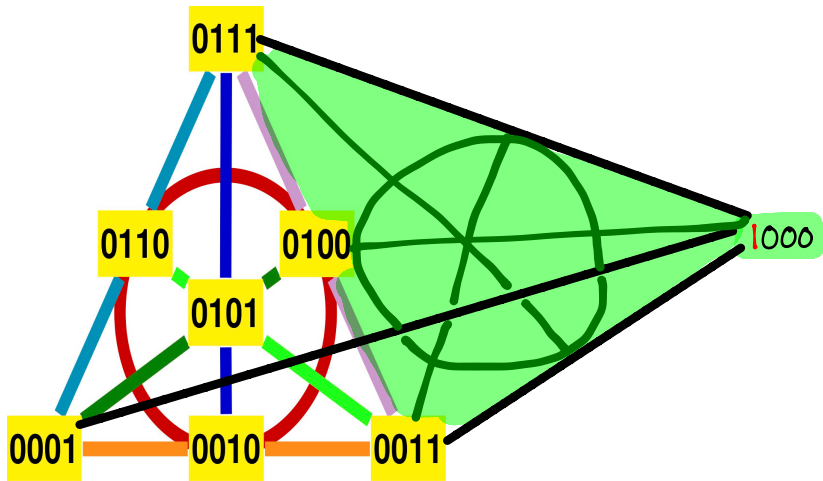
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Projective Geometry $PG(4)$

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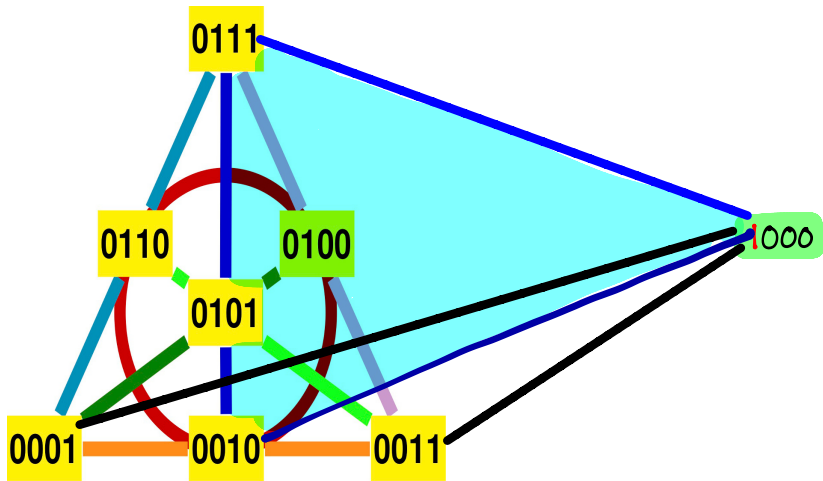
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Projective Geometry $PG(4)$

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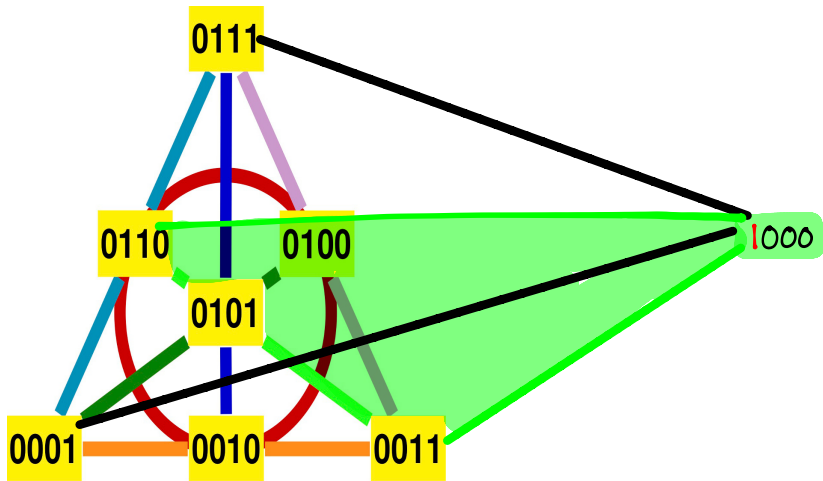
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Projective Geometry $PG(4)$

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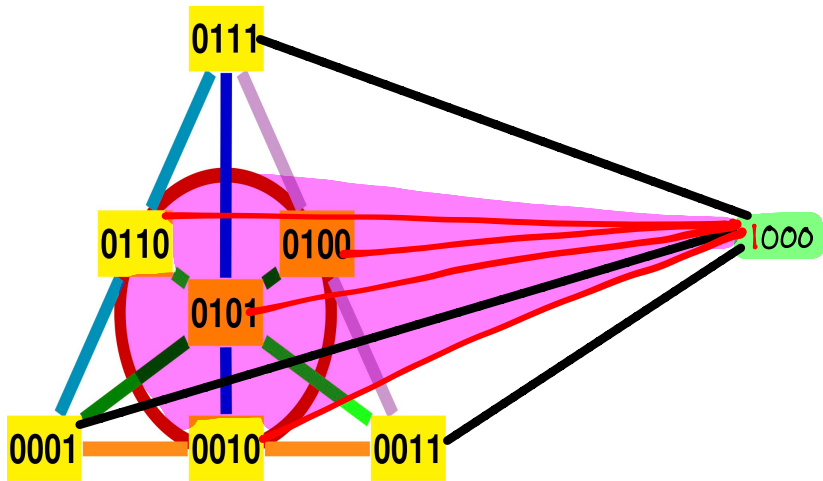
Lines: $\ell + m \equiv n \pmod{2}$



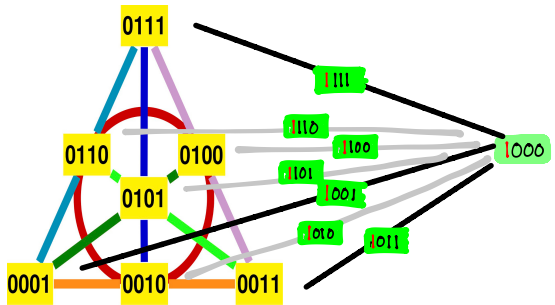
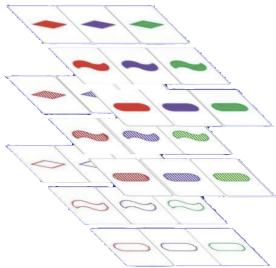
Projective Geometry $PG(4)$

Points: Nonzero binary points (p_1, p_2, p_3, p_4)

Lines: $\ell + m \equiv n \pmod{2}$

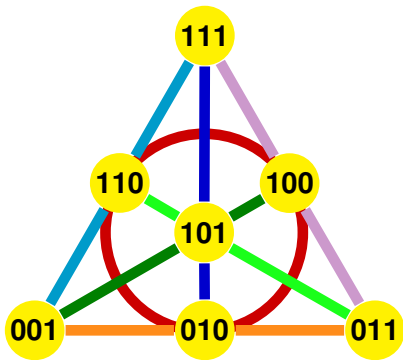


Projective Geometries “grow” much like Affine Geometries



Caps

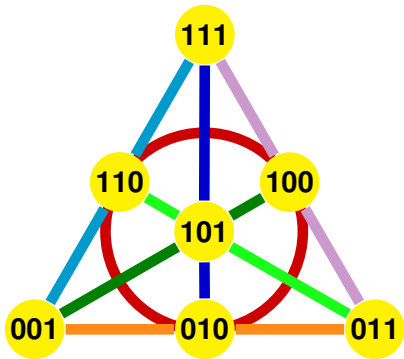
What's the *one* special feature of \mathcal{X} 's and \mathcal{O} 's points?



Caps

What's the *one* special feature of \mathcal{X} 's and \mathcal{O} 's points?

Cap: A set of points that contains no line.

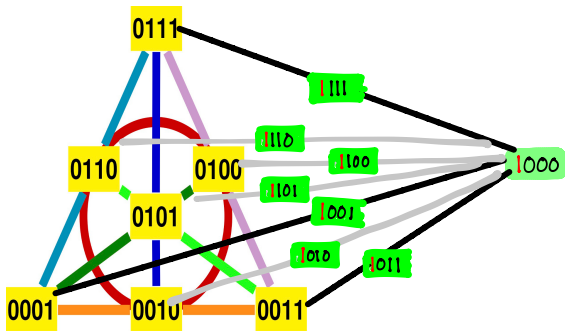


$PG(3)$ Max cap:
Rest is

Caps

What's the *one* special feature of \mathcal{X} 's and \mathcal{O} 's points?

Cap: A set of points that contains no line.



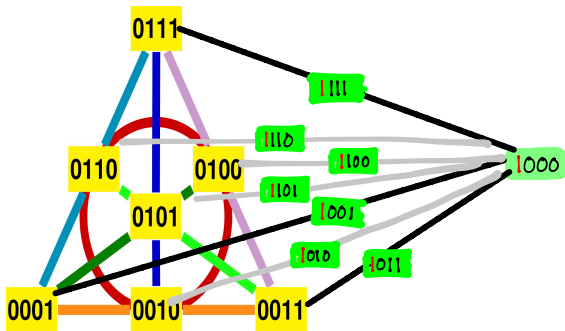
$PG(3)$ Max cap: 2^2
Rest is $PG(2)$

$PG(4)$ Max cap:
Rest is

Caps

What's the *one* special feature of \mathcal{X} 's and \mathcal{O} 's points?

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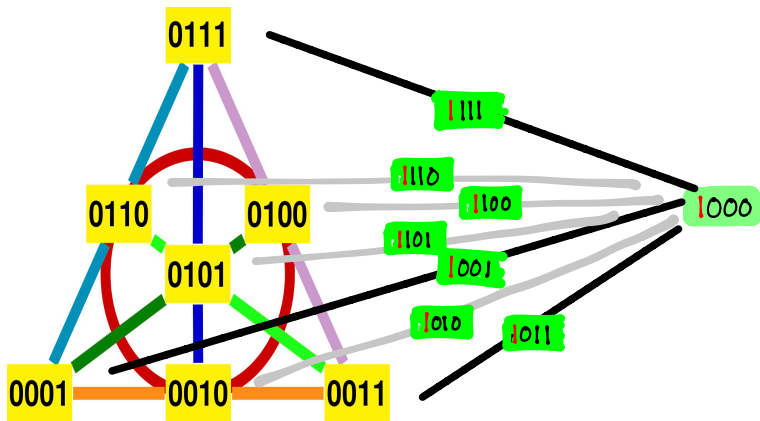


$PG(3)$ Max cap: 2^2
Rest is $PG(2)$

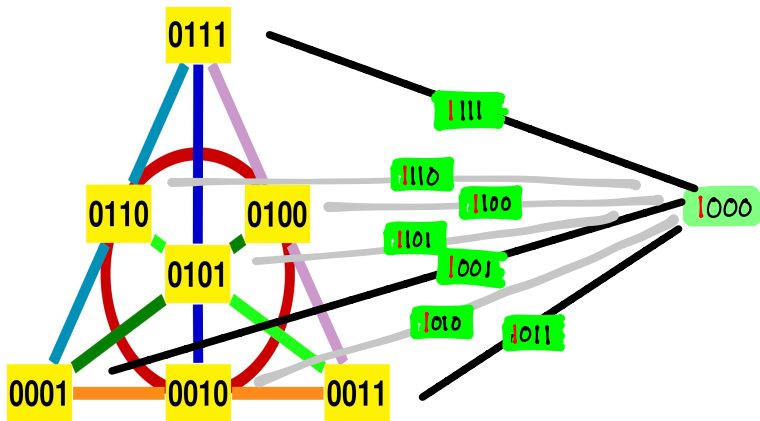
$PG(4)$ Max cap: 2^3
Rest is $PG(3)$

$PG(n)$ Max cap:
Rest is

How could we win Anti-Tic-Tac-Toe on $PG(n)$?

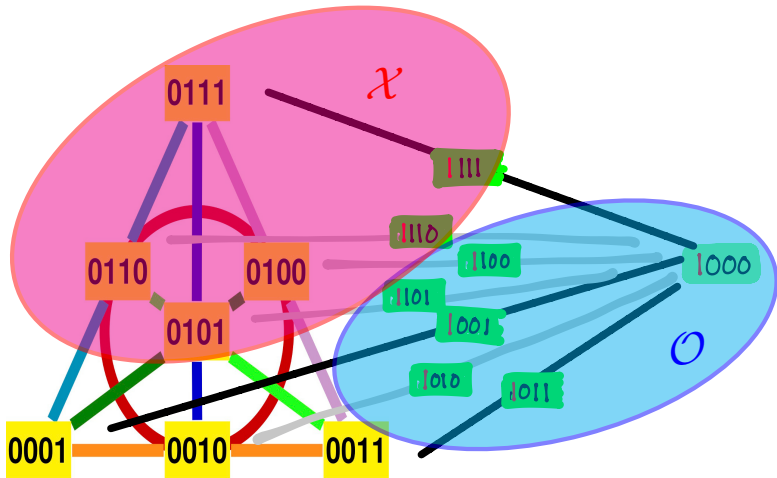


How could we win Anti-Tic-Tac-Toe on $PG(n)$?



Each player could (possibly) take: $\frac{3}{4}2^{n-1}$ ($\frac{3}{4}$ of a maximum cap!)

How could we win Anti-Tic-Tac-Toe on $PG(n)$?



Who wins?

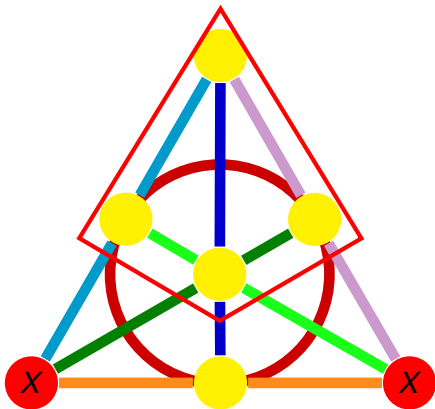
How to win Anti-Tic-Tac-Toe on $PG(n)$

Winning condition for \mathcal{O}

If \mathcal{O} can take $\frac{3}{4}$ of the points in a max cap,
then \mathcal{X} must lose the game on their next move.

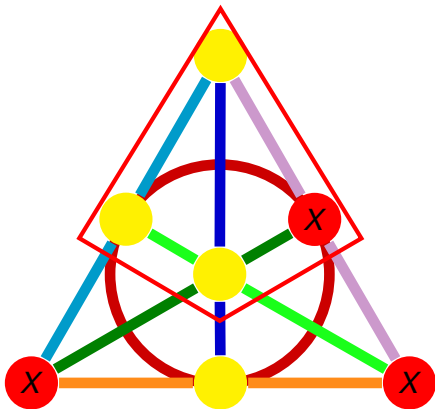
Pasch tiles

A maximum cap can be *tiled*:



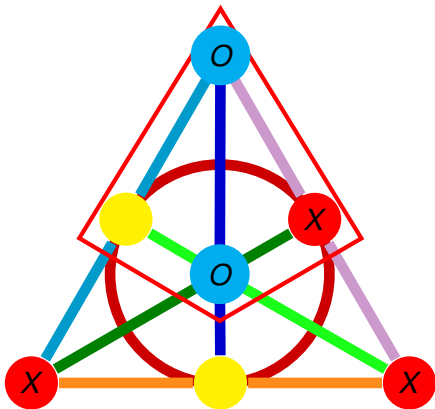
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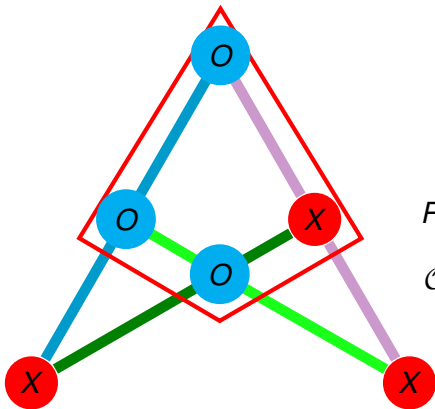
Pasch tiles

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Pasch tiles

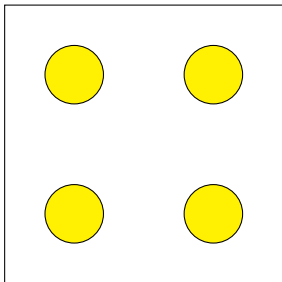
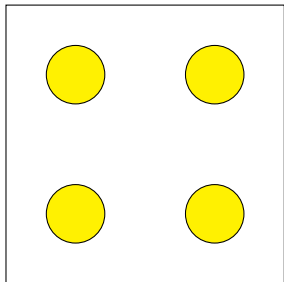
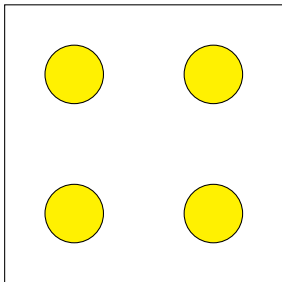
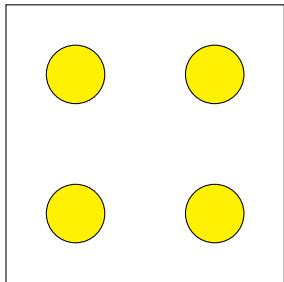
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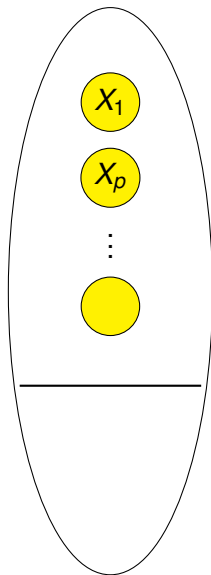
The
Pasch Configuration

O can take $\frac{3}{4}$ points
in each tile

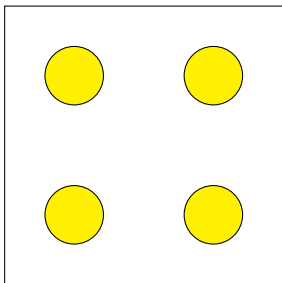
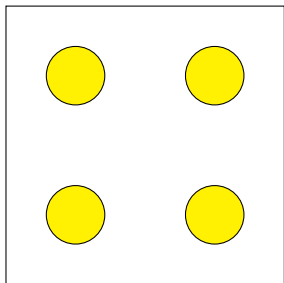
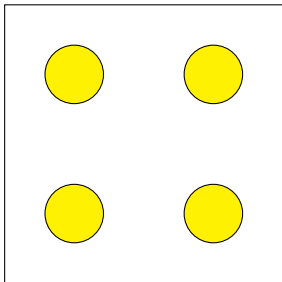
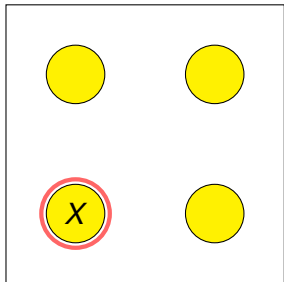
O's Pasch tiles



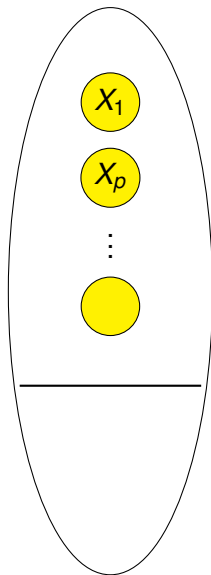
Rest of PG



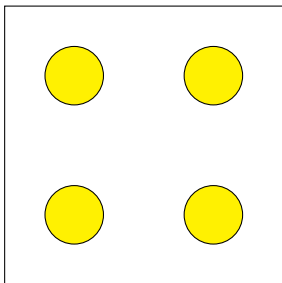
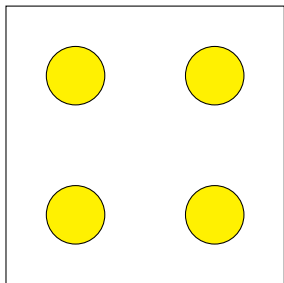
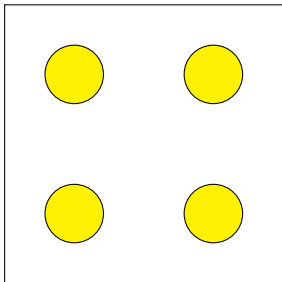
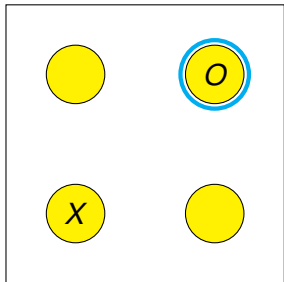
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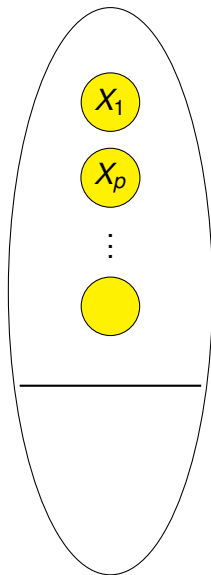
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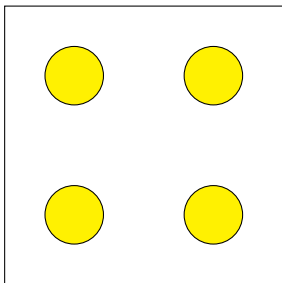
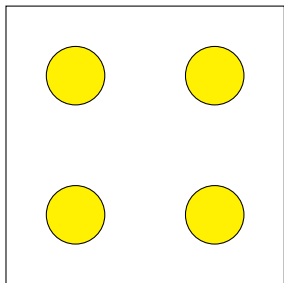
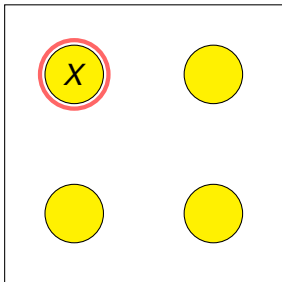
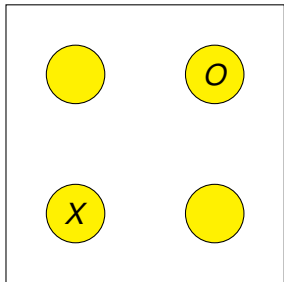
O 's Pasch tiles



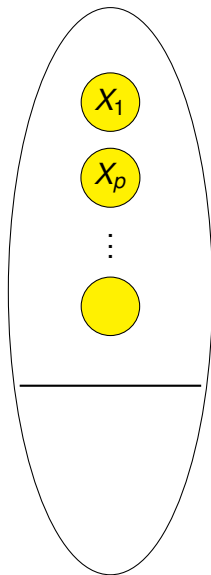
Rest of PG



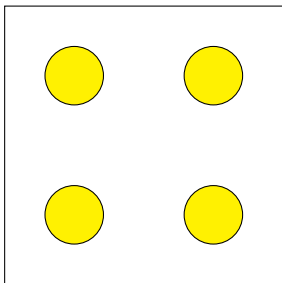
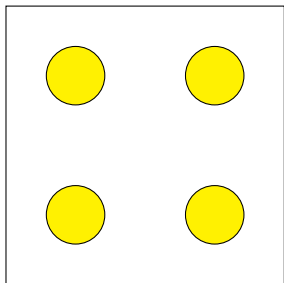
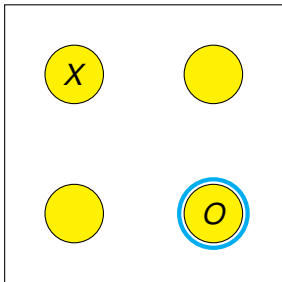
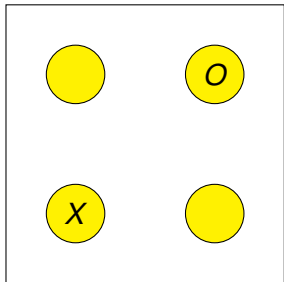
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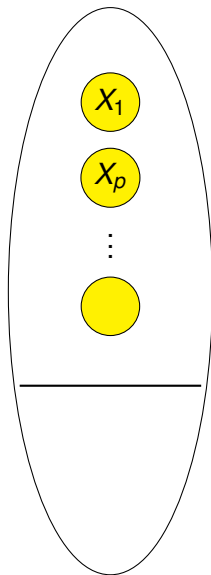
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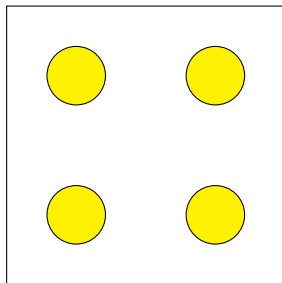
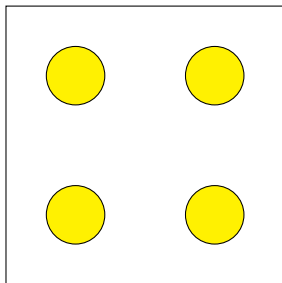
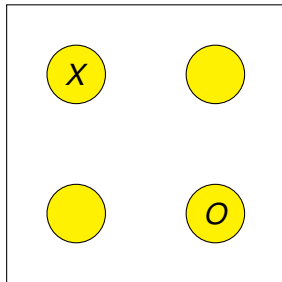
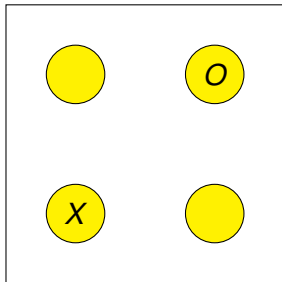
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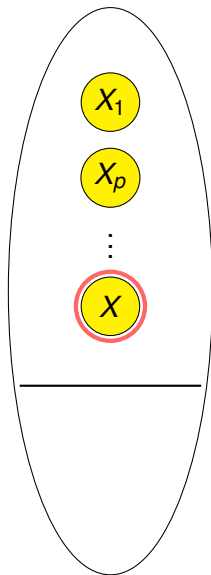
Rest of PG



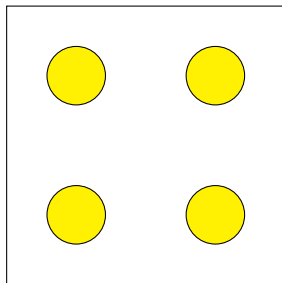
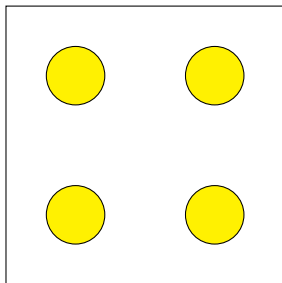
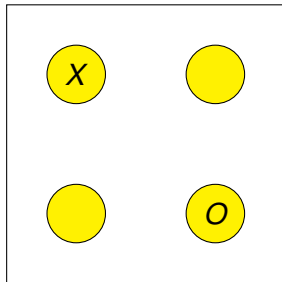
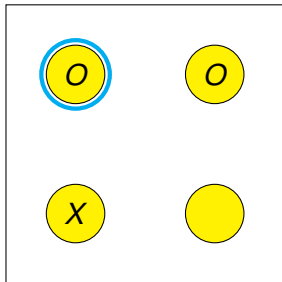
O's Pasch tiles



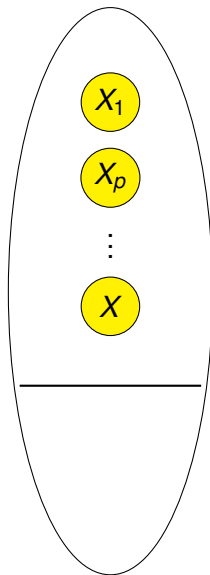
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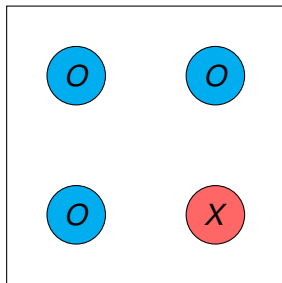
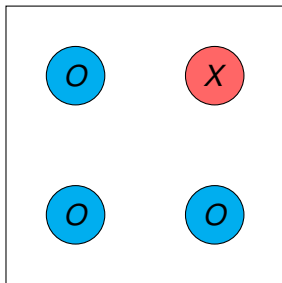
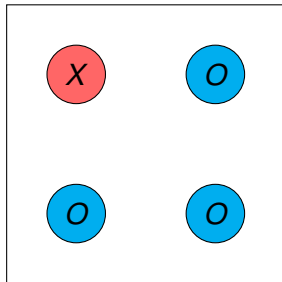
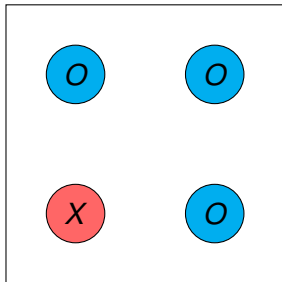
O's Pasch tiles



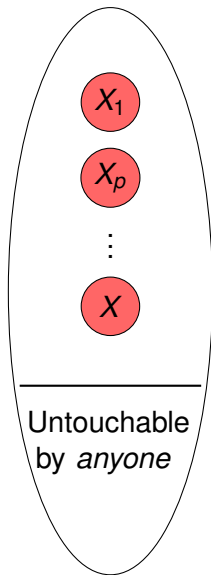
Rest of PG



O's Pasch tiles



Rest of PG



Summary

Anti-tic-tac-toe on Steiner triple systems

- SET = Affine Geometry: \mathcal{X} wins! (Using mitres)
- Projective Geometry: \mathcal{O} wins! (using Pasch configurations)
- What else is there to do?

Questions?

